

The polyhedral structure and complexity of multistage stochastic linear problem with general cost distribution

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Multistage stochastic linear programming (MSLP)

$$\begin{aligned} \min_{(\mathbf{x}_t)_{t \in [t_{\max}]}} \quad & \mathbb{E} \left[\sum_{t=1}^{t_{\max}} \mathbf{c}_t^\top \mathbf{x}_t \right] \\ \text{s.t.} \quad & \mathbf{T}_t \mathbf{x}_{t-1} + \mathbf{W}_t \mathbf{x}_t \leq \mathbf{h}_t & \forall t \in [t_{\max}] \\ & \mathbf{x}_t \text{ random variable in } \mathbb{R}^{n_t} & \forall t \in [t_{\max}] \\ & \mathbf{x}_t \in \sigma(\mathbf{c}_k, \mathbf{T}_k, \mathbf{W}_k, \mathbf{h}_k)_{k \leq t} & \forall t \in [t_{\max}] \\ & \mathbf{x}_0 \equiv x_0 \text{ given} \end{aligned}$$

where $\mathbf{c}_t \in \mathbb{R}^{n_t}$, $\mathbf{T}_t \in \mathbb{R}^{q_t \times n_{t-1}}$, $\mathbf{W}_t \in \mathbb{R}^{q_t \times n_t}$ and $\mathbf{h}_t \in \mathbb{R}^{q_t}$ are given random variables.

$(\mathbf{c}_t, \mathbf{T}_t, \mathbf{W}_t, \mathbf{h}_t)_{t \in [t_{\max}]}$ is an independent sequence.

We set $V_{t_{\max}+1} \equiv 0$ and:

$$V_t(x_{t-1}) := \mathbb{E} \left[\begin{array}{l} \min_{\mathbf{x}_t \in \mathbb{R}^{n_t}} \mathbf{c}_t^\top \mathbf{x}_t + V_{t+1}(x_t) \\ \text{s.t. } \mathbf{T}_t x_{t-1} + \mathbf{W}_t \mathbf{x}_t \leq \mathbf{h}_t \end{array} \right]$$

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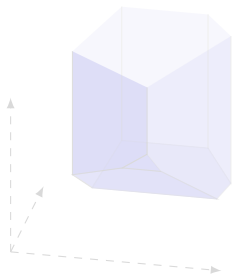
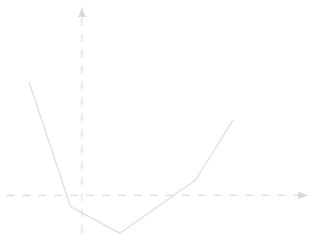
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Is V polyhedral ?

$$V(x) = \mathbb{E} \left[\min_{y \in \mathbb{R}^m} \begin{array}{l} \mathbf{c}^\top y + R(y) \\ \text{s.t. } \mathbf{T}x + \mathbf{W}y \leq \mathbf{h} \end{array} \right] = \mathbb{E} \left[\min_{y \in \mathbb{R}^m} (\mathbf{c}^\top y + R(y) + \mathbb{I}_{\mathbf{T}x + \mathbf{W}y \leq \mathbf{h}}) \right]$$

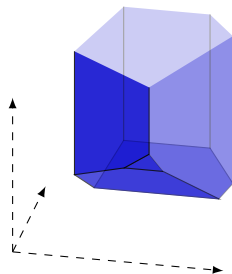
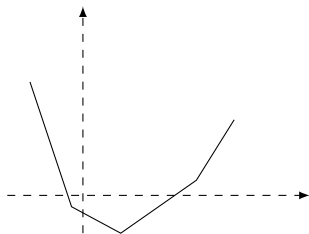
Question : On which conditions on the random variable \mathbf{c} , \mathbf{T} , \mathbf{W} and \mathbf{h} , is V **polyhedral** ?



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We can assume $R \equiv 0$

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If R is polyhedral, $\text{epi}(R) := \{(y, z) \mid Ay + b \leq z, Cy \leq d\}$

\rightsquigarrow We may assume $R \equiv 0$ by setting

$$\tilde{y} = \begin{pmatrix} y \\ z \end{pmatrix}, \tilde{\mathbf{c}} = \begin{pmatrix} \mathbf{c} \\ 1 \end{pmatrix}, \tilde{\mathbf{T}} = \begin{pmatrix} \mathbf{T} \\ 0 \end{pmatrix}, \tilde{\mathbf{W}} = \begin{pmatrix} \mathbf{W} & 0 \\ A & -1 \\ C & 0 \end{pmatrix}, \tilde{\mathbf{h}} = \begin{pmatrix} \mathbf{h} \\ -b \\ d \end{pmatrix}$$

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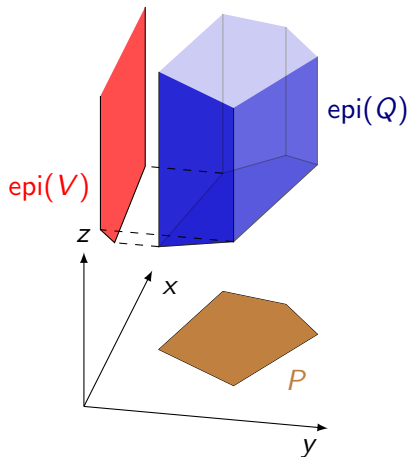
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c, T, W, h deterministic $\Rightarrow V$ polyhedral

For $x \in \mathbb{R}^n$,

$$\begin{aligned} V(x) &= \min_{y \in \mathbb{R}^m} (c^\top y + \mathbb{I}_{Tx+Wy \leq h}) \\ &= \min_{y \in \mathbb{R}^m} (c^\top y + \mathbb{I}_{(x,y) \in P}) \\ &= \min_{y \in \mathbb{R}^m} Q(x, y) \end{aligned}$$

V is polyhedral because
 $\text{epi}(V) \subset \mathbb{R}^{n+1}$ is the projection of
 $\text{epi}(Q) \subset \mathbb{R}^{n+m+1}$ on \mathbb{R}^{n+1} .



$\mathbf{c}, \mathbf{T}, \mathbf{W}, \mathbf{h}$ with finite support $\Rightarrow V$ polyhedral

Theorem (see e.g. Shapiro, Dentcheva, Ruszczyński)

If $\mathbf{c}, \mathbf{T}, \mathbf{W}, \mathbf{h}$ have a finite support, then V is polyhedral

Proof:

$$\begin{aligned} V(x) &= \sum_{k=1}^N p_k V_k(x) \\ &= \sum_{k=1}^N p_k \min_{y \in \mathbb{R}^m} (c_k^\top y + \mathbb{I}_{T_k x + W_k y \leq h_k}) \end{aligned}$$

where $p_k := \mathbb{P}[(\mathbf{c}, \mathbf{T}, \mathbf{W}, \mathbf{h}) = (c_k, T_k, W_k, h_k)]$.

Each V_k is polyhedral and $p_k \geq 0$.

\rightsquigarrow Question: are these assumptions tight ?

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Counter examples with stochastic constraints

Stochastic left hand
side constraint **T**

$$\begin{aligned} V(x) &= \mathbb{E} \left[\begin{array}{l} \min_{y \in \mathbb{R}^m} \quad y \\ \text{s.t.} \quad \mathbf{u}x \leq y \\ \quad \quad 1 \leq y \end{array} \right] \\ &= \mathbb{E} [\max(\mathbf{u}x, 1)] \\ &= \begin{cases} 1 & \text{if } x \leq 1 \\ \frac{x}{2} + \frac{1}{2x} & \text{if } x \geq 1 \end{cases} \end{aligned}$$

Stochastic right hand
side constraint **h**

$$\begin{aligned} V(x) &= \mathbb{E} \left[\begin{array}{l} \min_{y \in \mathbb{R}^m} \quad y \\ \text{s.t.} \quad \mathbf{u} \leq y \\ \quad \quad x \leq y \end{array} \right] \\ &= \mathbb{E} [\max(x, \mathbf{u})] \\ &= \begin{cases} \frac{1}{2} & \text{if } x \leq 0 \\ \frac{x^2+1}{2} & \text{if } x \in [0, 1] \\ x & \text{if } x \geq 1 \end{cases} \end{aligned}$$

where **u** is uniform on $[0, 1]$.

Remaining case: only \mathbf{c} stochastic

$$V(x) = \mathbb{E} \left[\min_{y \in \mathbb{R}^m} \mathbf{c}^\top y \quad \text{s.t.} \quad T\mathbf{x} + W\mathbf{y} \leq h \right] = \mathbb{E} \left[\min_{y \in \mathbb{R}^m} (\mathbf{c}^\top y + \mathbb{I}_{T\mathbf{x} + W\mathbf{y} \leq h}) \right]$$

Theorem (FGL 2020)

If T , W and h are deterministic, then for all distributions of \mathbf{c} such that V is well defined, V is polyhedral.

\leadsto This extends easily to \mathbf{T} , \mathbf{W} and \mathbf{h} with a finite support.

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Reformulation of $V(x)$ highlighting the role of the fiber P_x

For a given x ,

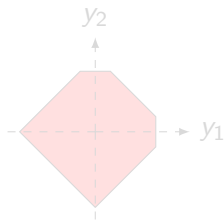
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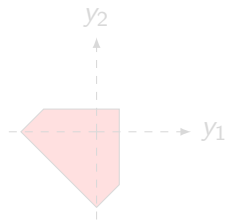
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Illustrative running example:

$$P_x := \{y \in \mathbb{R}^m \mid \|y\|_1 \leq 1, \quad y_1 \leq x, \quad y_2 \leq x\}$$



P_x for $x = 0.8$



P_x for $x = 0.3$

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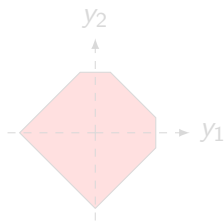
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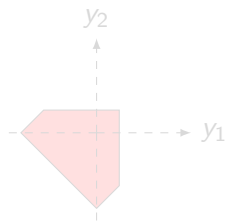
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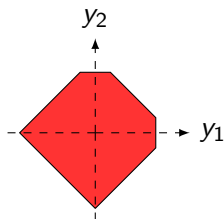
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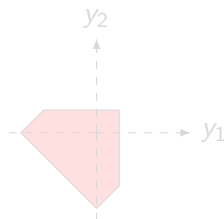
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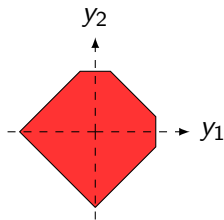
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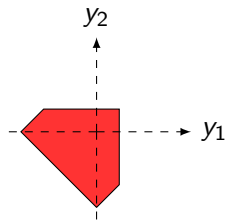
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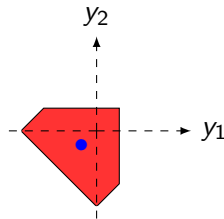
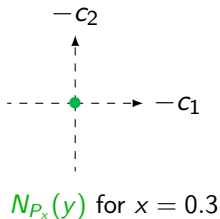
Normal fan $\mathcal{N}(P_x)$

Definition

The normal fan of the fiber P_x is

$$\mathcal{N}(P_x) := \{N_{P_x}(y) \mid y \in P_x\}$$

with $N_{P_x}(y) = \{c \mid \forall y' \in P_x, c^\top(y' - y) \leq 0\}$ the normal cone of P_x on y .



P_x and $N_{P_x}(y)$ for $x = 0.3$

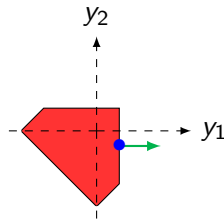
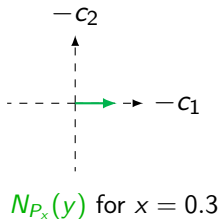
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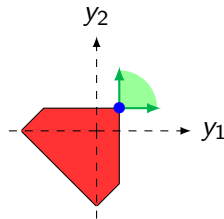
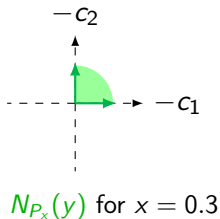
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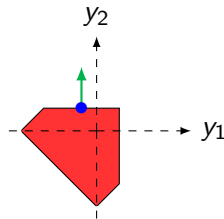
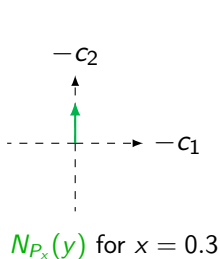
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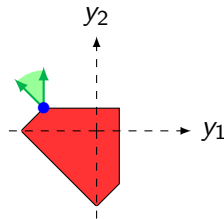
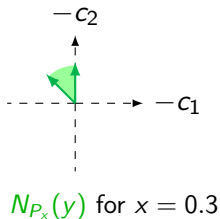
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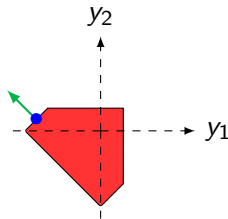
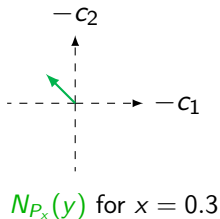
Normal fan $\mathcal{N}(P_x)$

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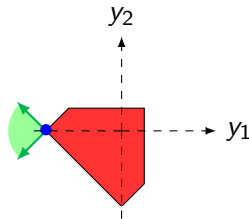
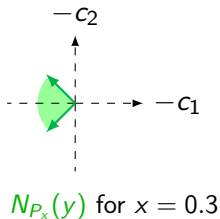
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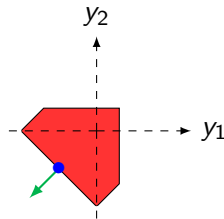
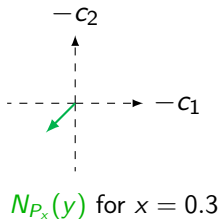
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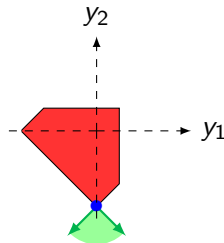
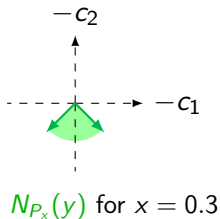
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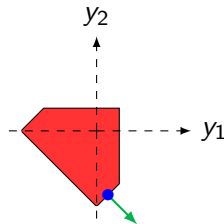
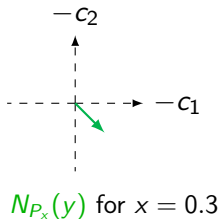
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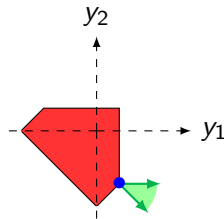
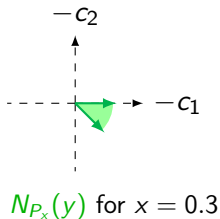
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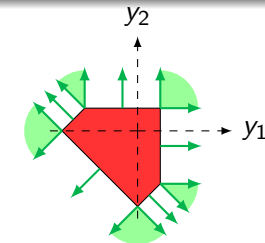
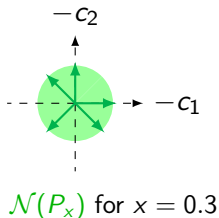
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Proposition

If P_x is bounded, $\{\text{ri}(N) \mid N \in \mathcal{N}(P_x)\}$ is a partition of \mathbb{R}^m .



P_x and $\mathcal{N}(P_x)$ for $x = 0.3$

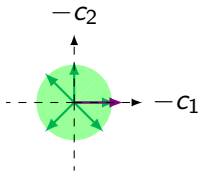
$\mathcal{N}(P_x)$: partition of $-\mathbf{c}$ coherent with the min

For a given x , we have

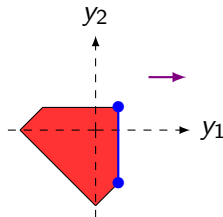
$$V(x) = \mathbb{E} \left[\min_{y \in P_x} \mathbf{c}^\top y \right]$$

For any $N \in \mathcal{N}(P_x)$ and $-\mathbf{c} \rightarrow \arg \min_{y \in P_x} \mathbf{c}^\top y$ is constant for all $-\mathbf{c} \in \text{ri}(N)$.

$\arg \min_{y \in P_x} \mathbf{c}^\top y$ is a face of P_x .



Cost $-\mathbf{c}$ and $\mathcal{N}(P_x)$ for $x = 0.3$



P_x for $x = 0.3$

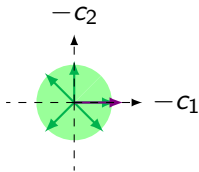
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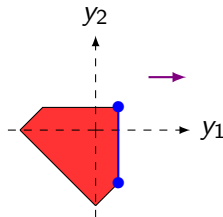
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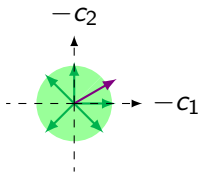
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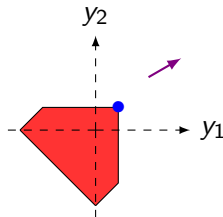
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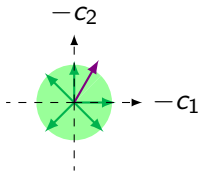
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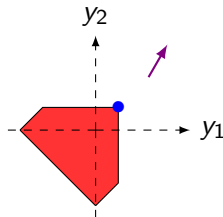
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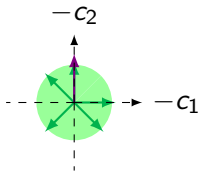
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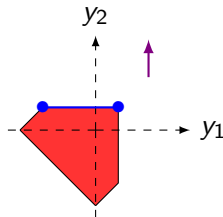
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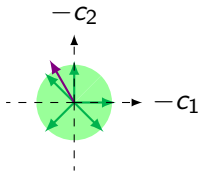
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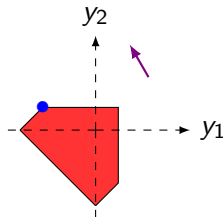
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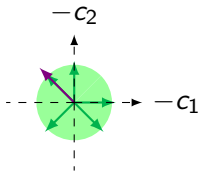
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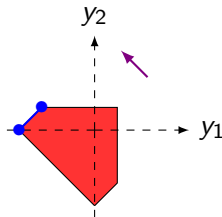
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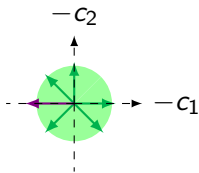
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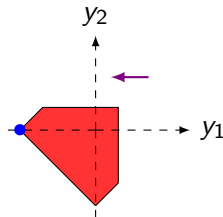
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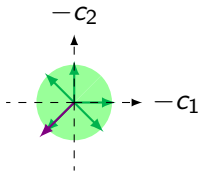
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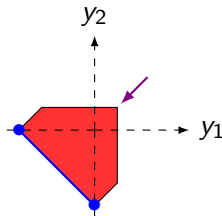
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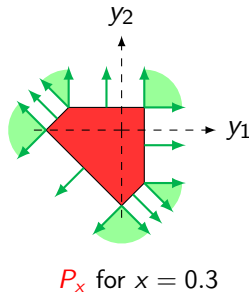
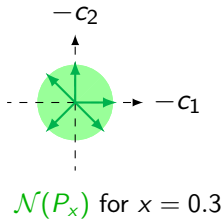
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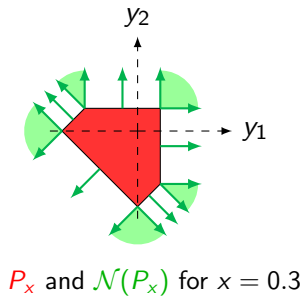
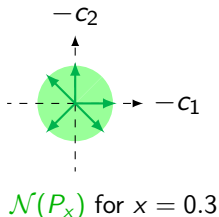


Reduction to a finite sum

For a fixed x ,

$$V(x) = \mathbb{E} \left[\min_{y \in P_x} \mathbf{c}^\top y \right] = \sum_{N \in \mathcal{N}(P_x)} \mathbb{E} \left[\mathbf{c}^\top \mathbf{1}_{\mathbf{c} \in -\text{ri}(N)} \right] y_N(x)$$

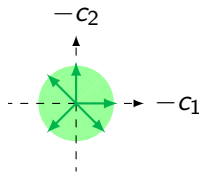
where $y_N(x) \in \arg \min_{y \in P_x} \mathbf{c}^\top y$ for any $\mathbf{c} \in \text{ri}(N)$.



General cost \mathbf{c} is equivalent to discrete cost $\check{\mathbf{c}}$ for given x

For a fixed x ,

$$\begin{aligned} V(x) &= \mathbb{E} \left[\min_{y \in P_x} \mathbf{c}^\top y \right] \\ &= \sum_{N \in \mathcal{N}(P_x)} \mathbb{E} \left[\mathbf{c}^\top \mathbf{1}_{\mathbf{c} \in -\text{ri } N} \right] y_N(x) \end{aligned}$$



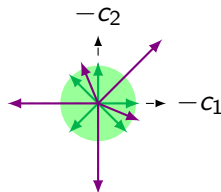
$\mathcal{N}(P_x)$ for $x = 0.3$

We draw a continuous cost \mathbf{c} .

General cost \mathbf{c} is equivalent to discrete cost $\check{\mathbf{c}}$ for given x

For a fixed x ,

$$\begin{aligned} V(x) &= \mathbb{E} \left[\min_{y \in P_x} \mathbf{c}^\top y \right] \\ &= \sum_{N \in \mathcal{N}(P_x)} \mathbb{E} [\mathbf{c}^\top \mathbf{1}_{\mathbf{c} \in -\text{ri } N}] y_N(x) \\ &= \sum_{N \in \mathcal{N}(P_x)} p_N \check{\mathbf{c}}_N y_N(x) \end{aligned}$$



$\mathcal{N}(P_x)$ and $p_N \check{\mathbf{c}}_N$ for $x = 0.3$

For $N \in \mathcal{N}(P_x)$,

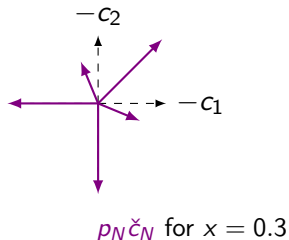
$$\begin{aligned} p_N &:= \mathbb{P} [\mathbf{c} \in -\text{ri } N] \\ \check{\mathbf{c}}_N &:= \mathbb{E} [\mathbf{c} | \mathbf{c} \in -\text{ri } N] \end{aligned}$$

Instead of drawing a general \mathbf{c} ,
we draw a discrete cost $\check{\mathbf{c}}$ indexed by
the finite collection $\mathcal{N}(P_x)$.

General cost \mathbf{c} is equivalent to discrete cost $\check{\mathbf{c}}$ for given x

For a fixed x ,

$$\begin{aligned}
 V(x) &= \mathbb{E} \left[\min_{y \in P_x} \mathbf{c}^\top y \right] \\
 &= \sum_{N \in \mathcal{N}(P_x)} \mathbb{E} \left[\mathbf{c}^\top \mathbf{1}_{\mathbf{c} \in -ri N} \right] y_N(x) \\
 &= \sum_{N \in \mathcal{N}(P_x)} p_N \check{c}_N y_N(x) \\
 &= \sum_{N \in \mathcal{N}(P_x)} p_N \min_{y \in P_x} \check{c}_N^\top y
 \end{aligned}$$



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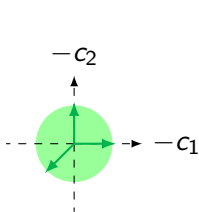
Contents

- 1 Introduction
- 2 Studying the polyhedral structure of cost-to-go functions
 - Fixed state x and normal fan
 - Variable state x and chamber complex
 - Main theorem
- 3 Computation and formulas
- 4 Complexity results

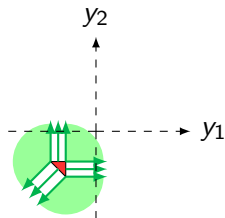
$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

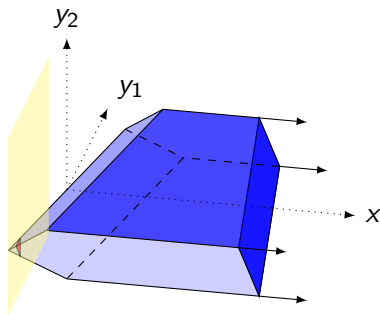
$$x = -0.4$$



$\mathcal{N}(P_x)$



P_x and $\mathcal{N}(P_x)$



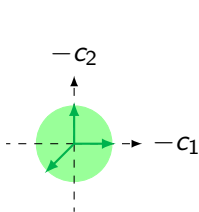
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P and P_x

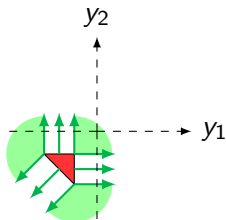
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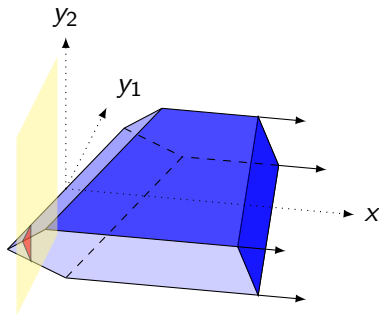
$$x = -0.3$$



$\mathcal{N}(P_x)$



P_x and $\mathcal{N}(P_x)$



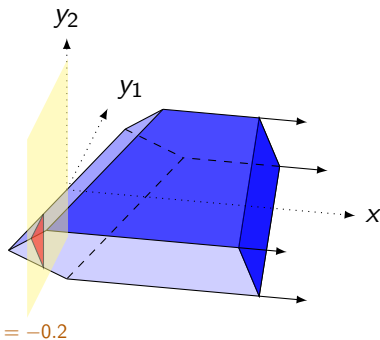
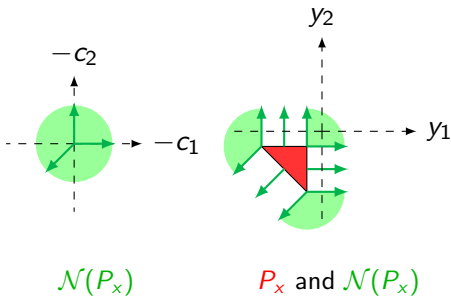
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P and P_x

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$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

$$x = -0.2$$

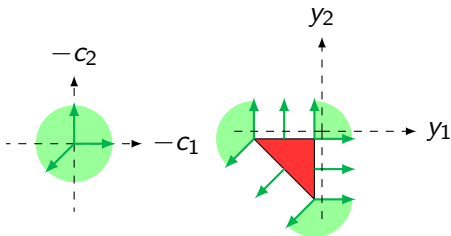


P and P_x

$\mathcal{N}(P_x)$ is piecewise constant with x .

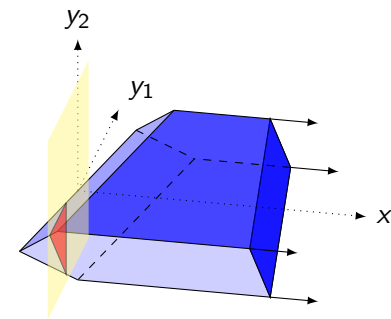
$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

$$x = -0.1$$



$\mathcal{N}(P_x)$

P_x and $\mathcal{N}(P_x)$



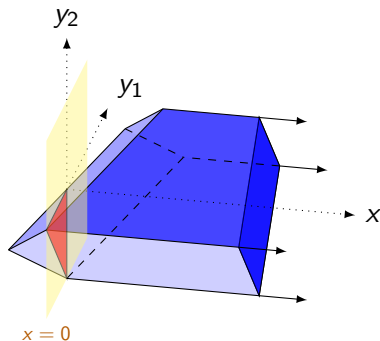
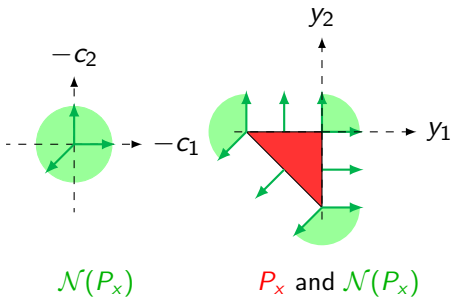
$$x = -0.1$$

P and P_x

$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

$$x = 0$$

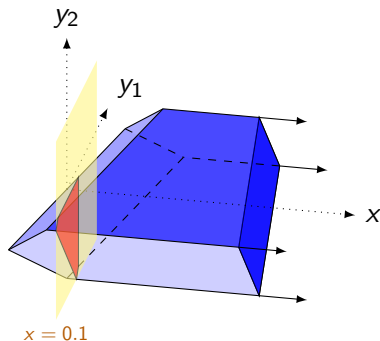
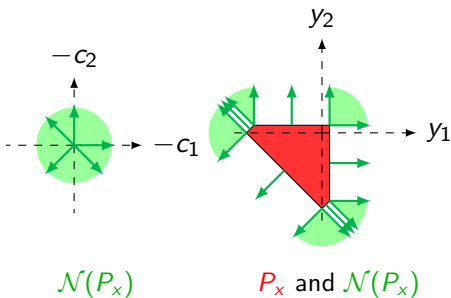


$$P \text{ and } P_x$$

$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

$$x = 0.1$$

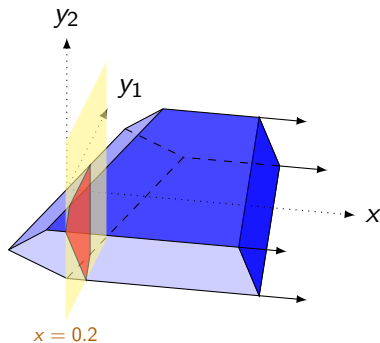
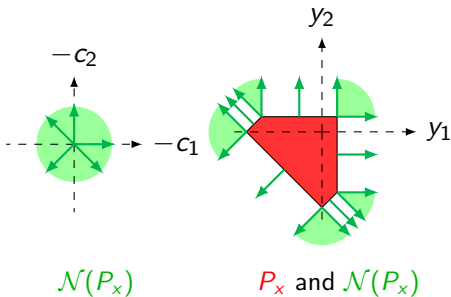


P and P_x

$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

$$x = 0.2$$

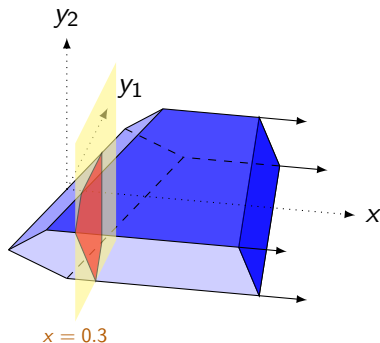
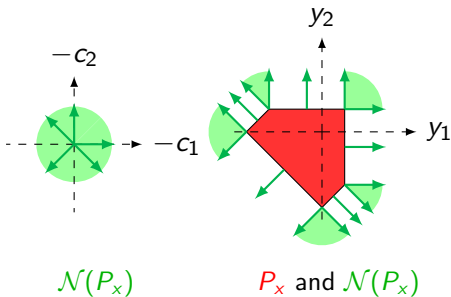


$$P \text{ and } P_x$$

$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

$$x = 0.3$$

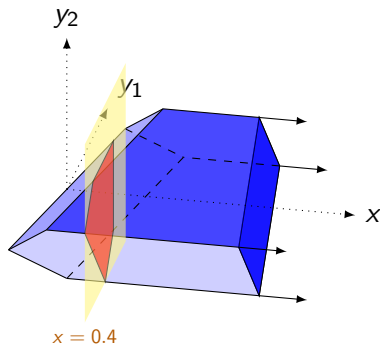
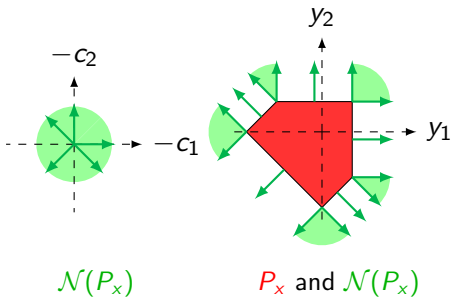


P and P_x

$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

$$x = 0.4$$

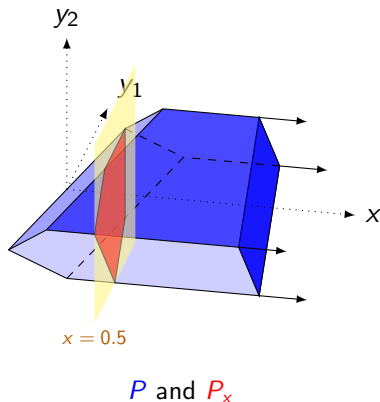
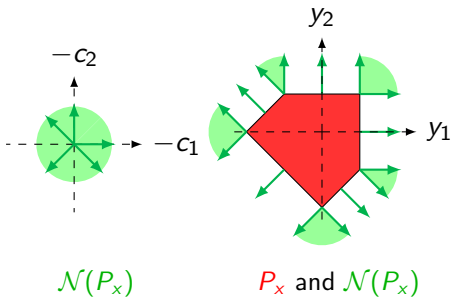


$$P \text{ and } P_x$$

$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

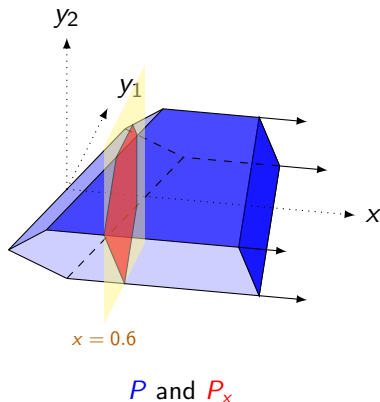
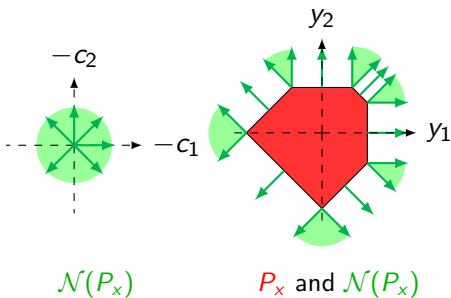
$$x = 0.5$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

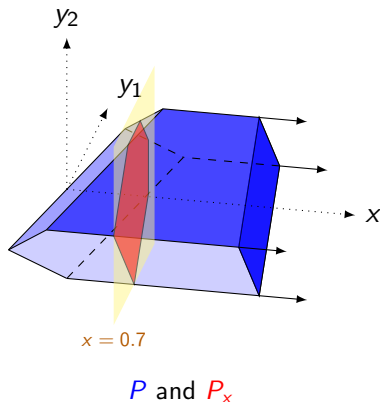
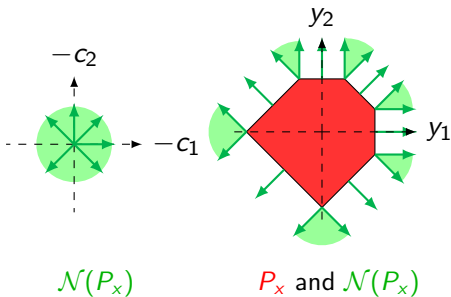
$$x = 0.6$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

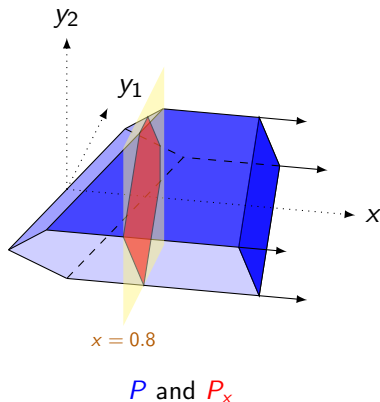
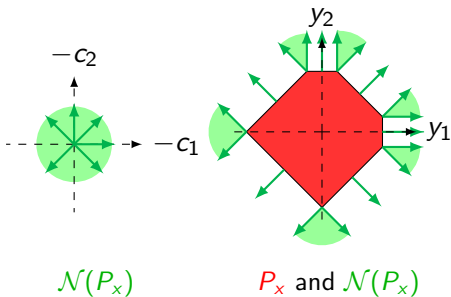
$$x = 0.7$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

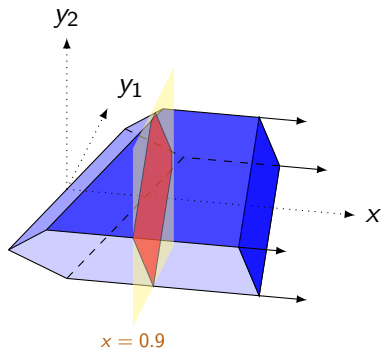
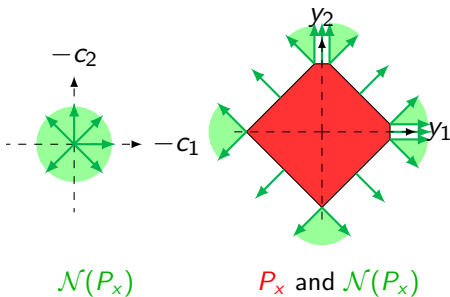
$$x = 0.8$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

$$x = 0.9$$

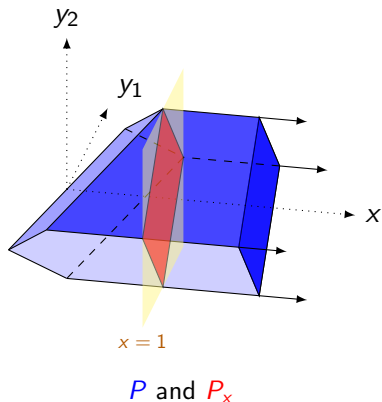
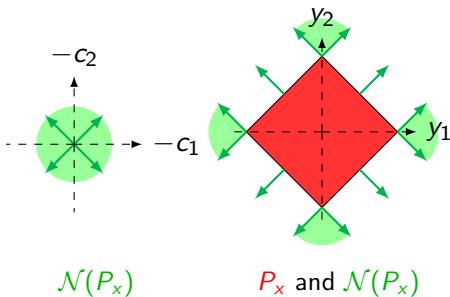


$$P \text{ and } P_x$$

$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

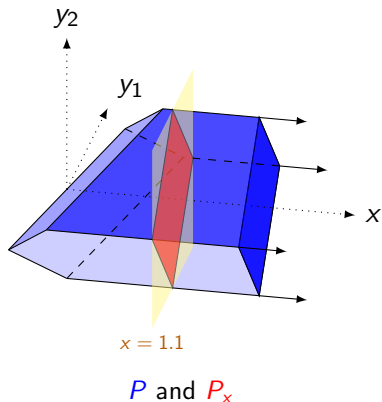
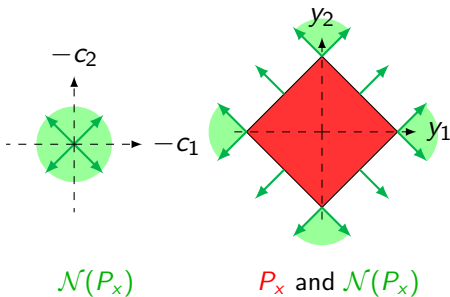
$$x = 1$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

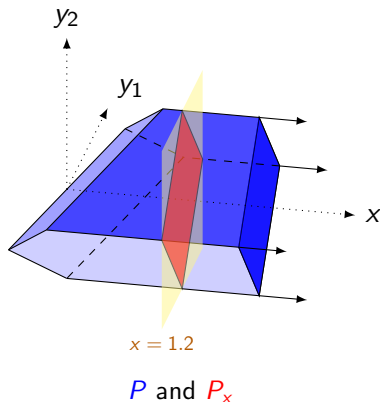
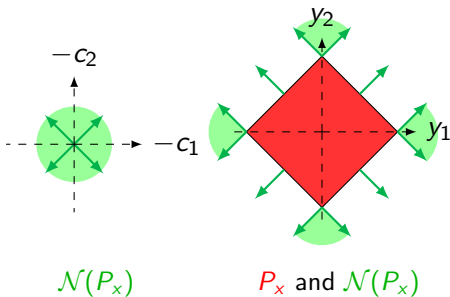
$$x = 1.1$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

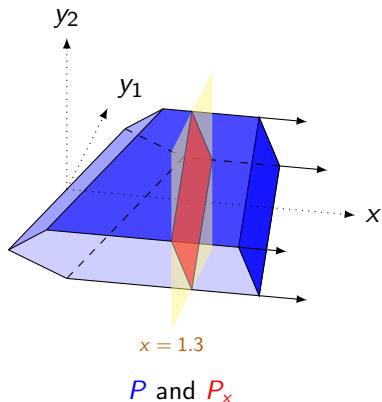
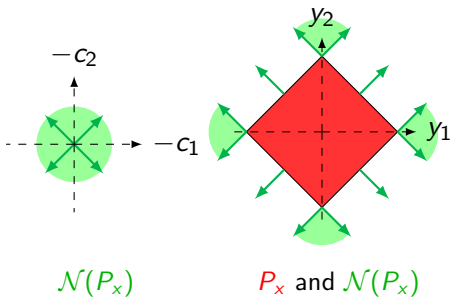
$$x = 1.2$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

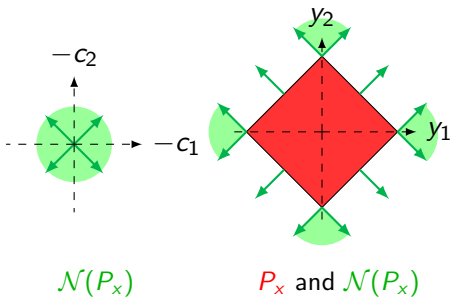
$$x = 1.3$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

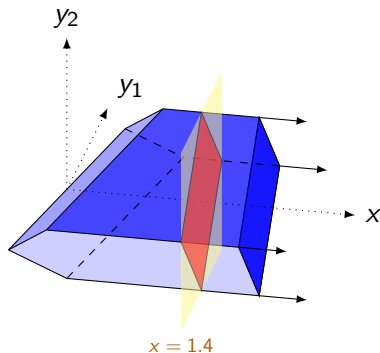
$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

$$x = 1.4$$



$\mathcal{N}(P_x)$

P_x and $\mathcal{N}(P_x)$

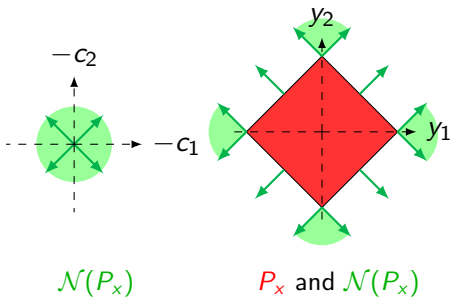


P and P_x

$\mathcal{N}(P_x)$ is piecewise constant with x .

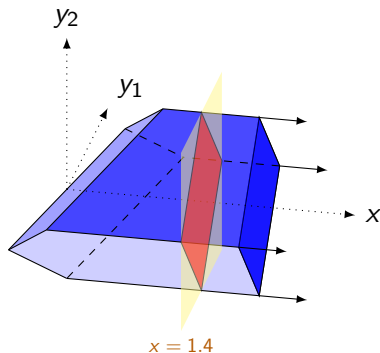
$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

$$x = 1.4$$



$\mathcal{N}(P_x)$

P_x and $\mathcal{N}(P_x)$

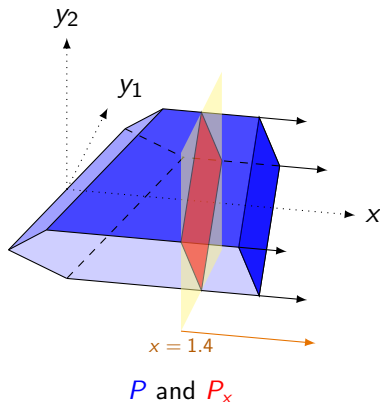
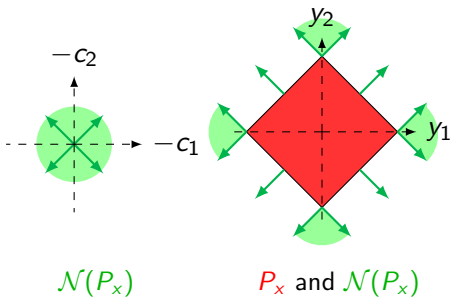


P and P_x

$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

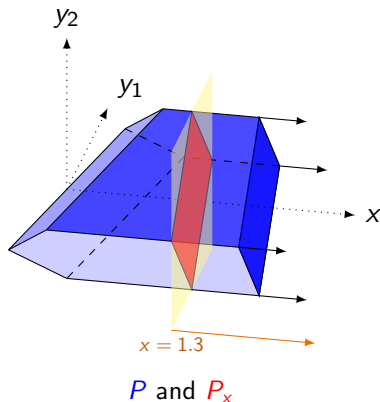
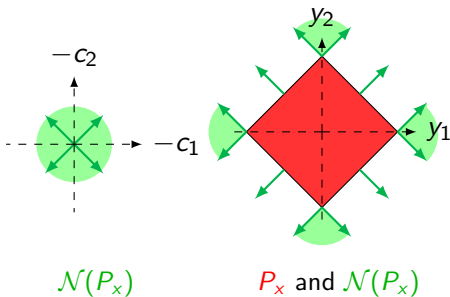
$$x = 1.4$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

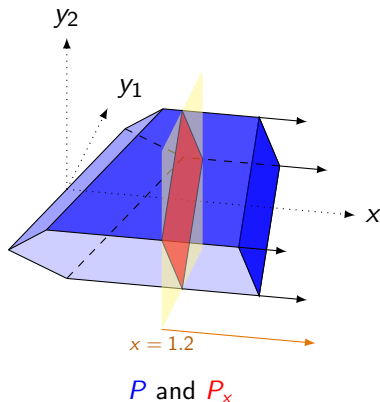
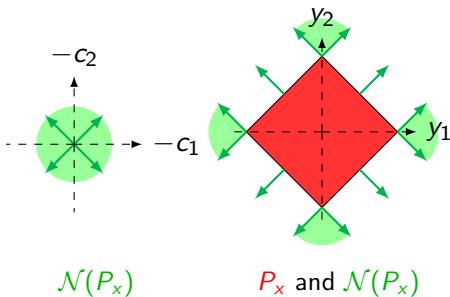
$$x = 1.3$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

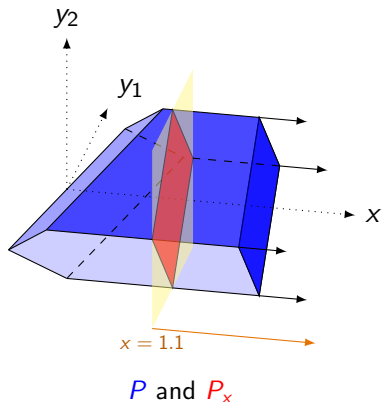
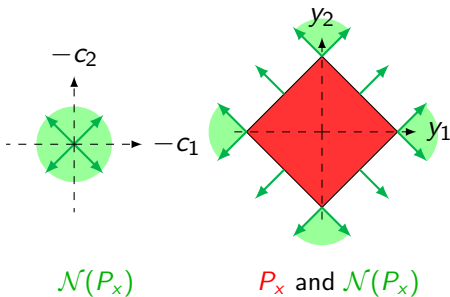
$$x = 1.2$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

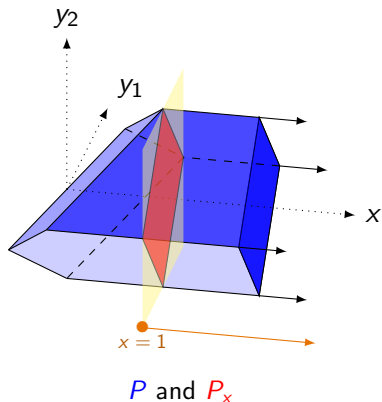
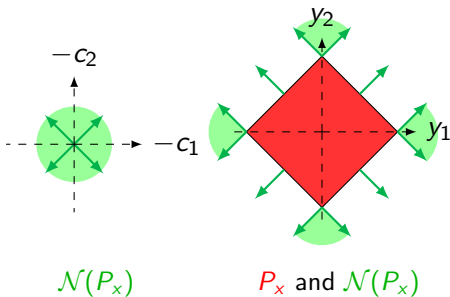
$$x = 1.1$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

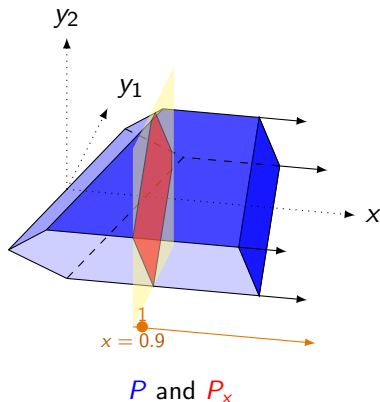
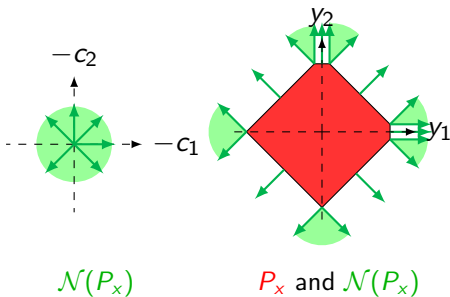
$$x = 1$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

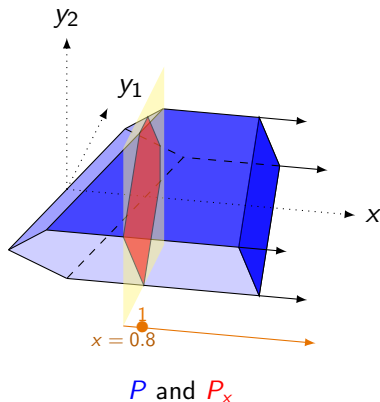
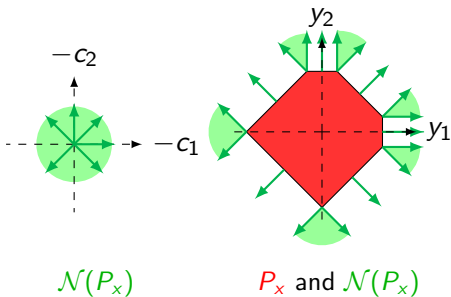
$$x = 0.9$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

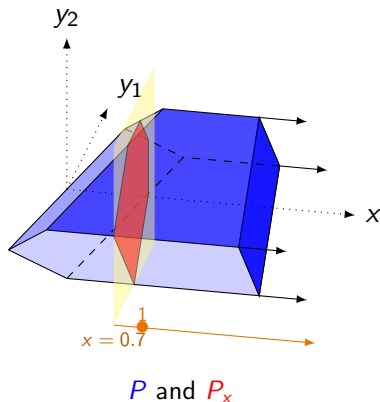
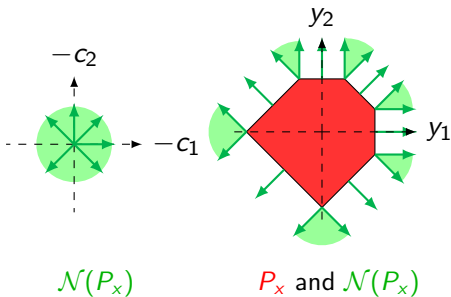
$$x = 0.8$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

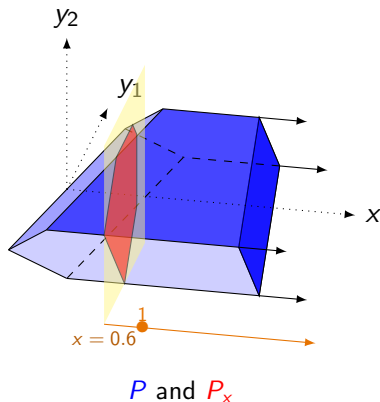
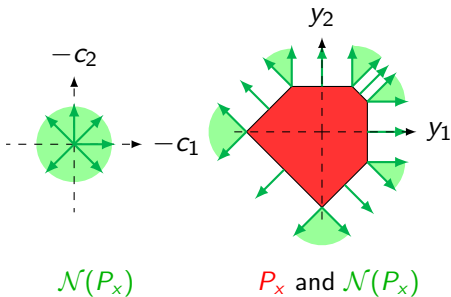
$$x = 0.7$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

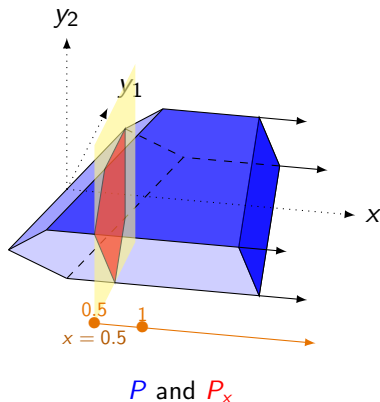
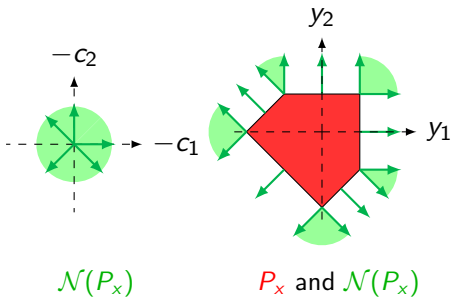
$$x = 0.6$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

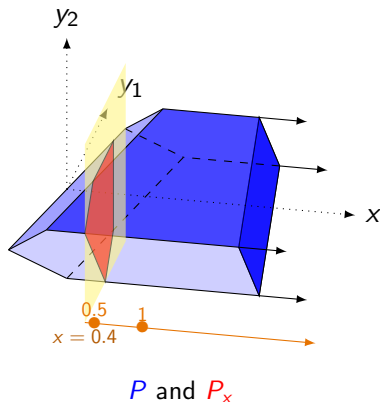
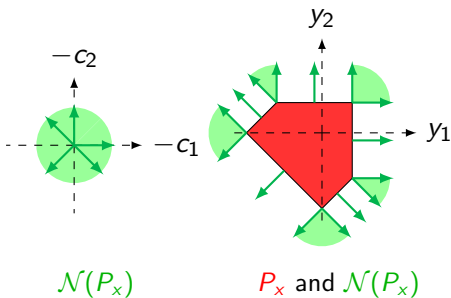
$$x = 0.5$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

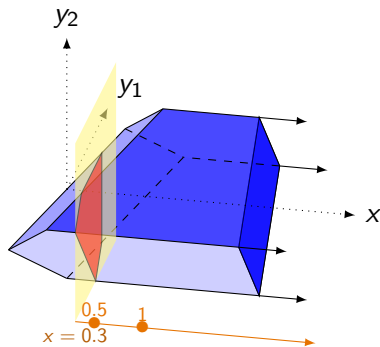
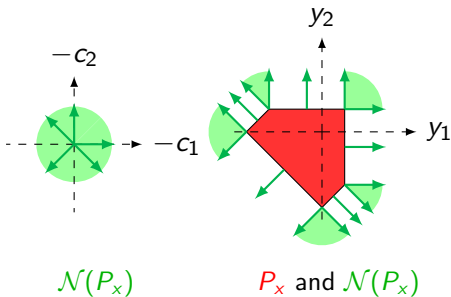
$x = 0.4$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

$$x = 0.3$$

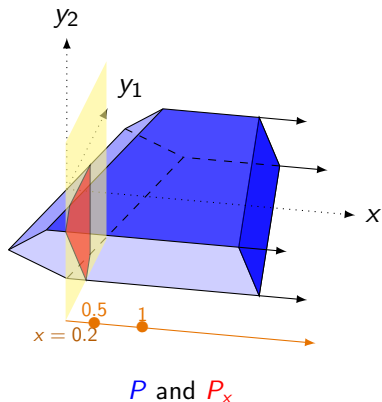
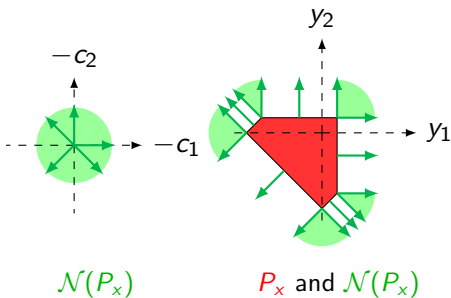


$$P \text{ and } P_x$$

$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

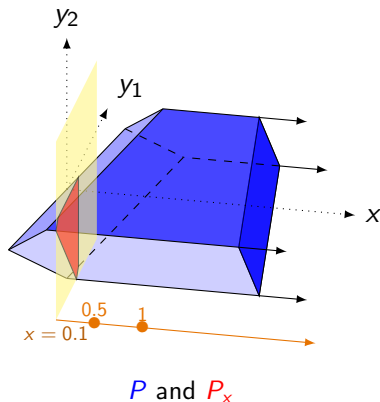
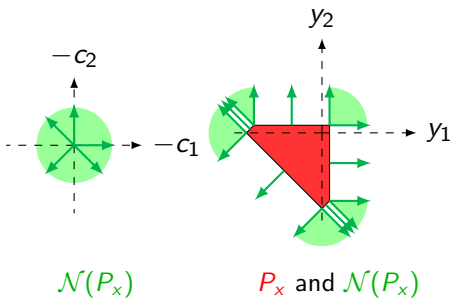
$$x = 0.2$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

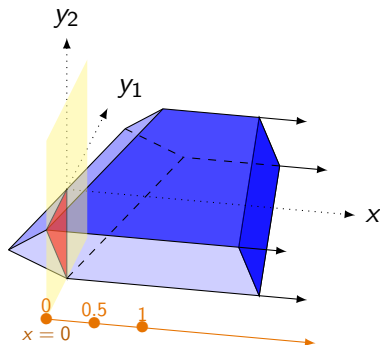
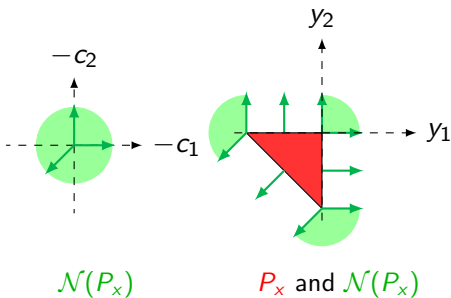
$$x = 0.1$$



$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

$$x = 0$$

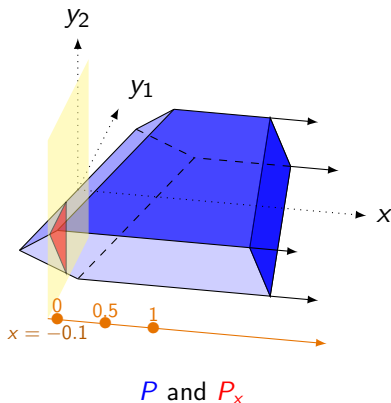
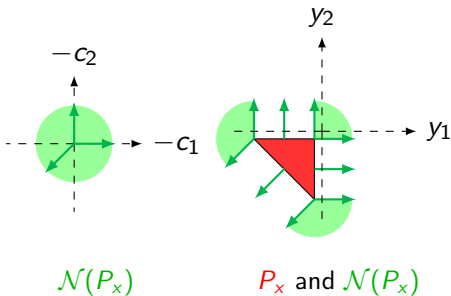


$$P \text{ and } P_x$$

$\mathcal{N}(P_x)$ is piecewise constant with x .

$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

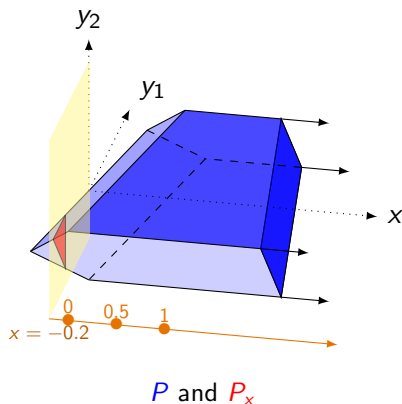
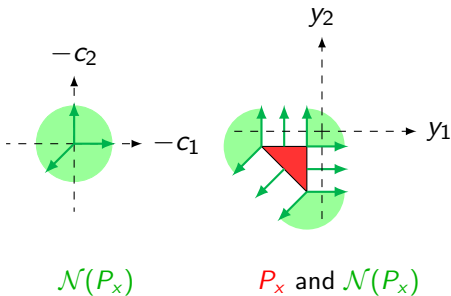
$$x = -0.1$$



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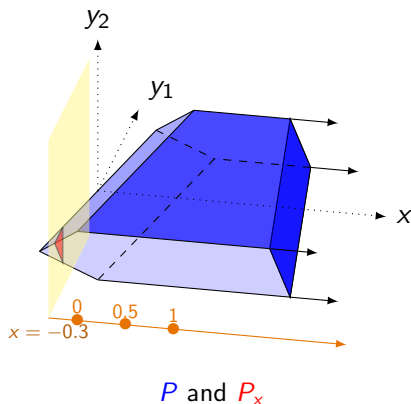
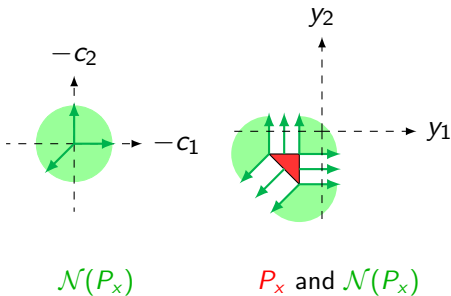
$$x = -0.2$$



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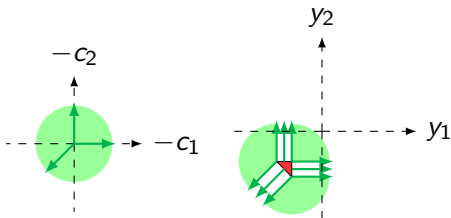
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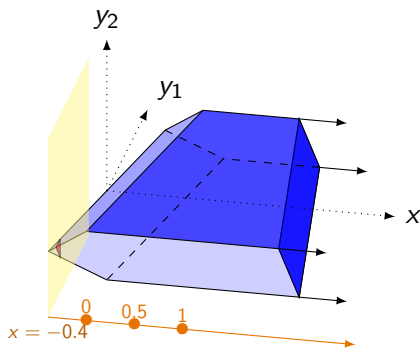
$$P := \{(x, y) \mid Tx + Wy \leq h\} \quad \text{and} \quad P_x := \{y \mid Tx + Wy \leq h\}$$

$$x = -0.4$$



$\mathcal{N}(P_x)$

P_x and $\mathcal{N}(P_x)$



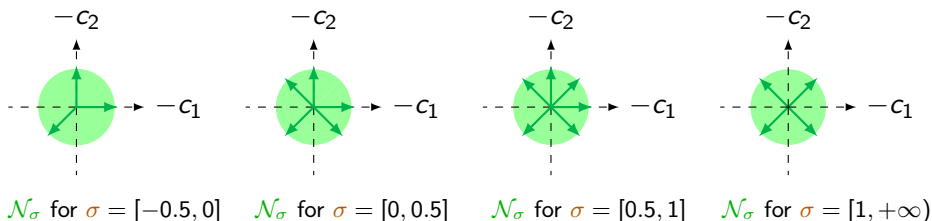
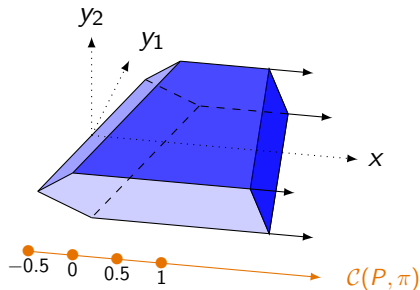
P and P_x

What are the constant regions of $\mathcal{N}(P_x)$?

Lemma

There exists a collection $\mathcal{C}(P, \pi)$ called the **chamber complex** whose relative interior of cells are the constant regions of $x \rightarrow \mathcal{N}(P_x)$.

For $\sigma \in \mathcal{C}(P, \pi)$ and $x, x' \in \text{ri}(\sigma)$,
 $\mathcal{N}(P_x) = \mathcal{N}(P_{x'}) =: \mathcal{N}_\sigma$



Chamber complex

Definition (Billera, Sturmfels 92)

The *chamber complex* $\mathcal{C}(P, \pi)$ of P along π is

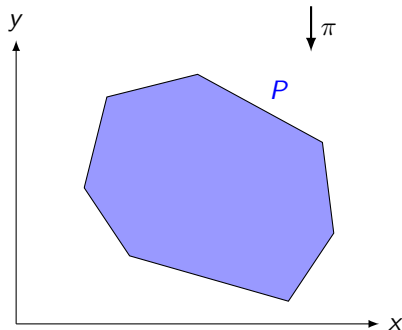
$$\mathcal{C}(P, \pi) := \{\sigma_{P, \pi}(x) \mid x \in \pi(P)\}$$

where

$$\sigma_{P, \pi}(x) := \bigcap_{F \in \mathcal{F}(P) \text{ s.t. } x \in \pi(F)} \pi(F)$$

where $\mathcal{F}(P)$ is the set of faces of P
and π is the projection $(x, y) \rightarrow x$

$$\pi(E) := \{x \in \mathbb{R}^n \mid \exists y \in \mathbb{R}^m, (x, y) \in E\}$$



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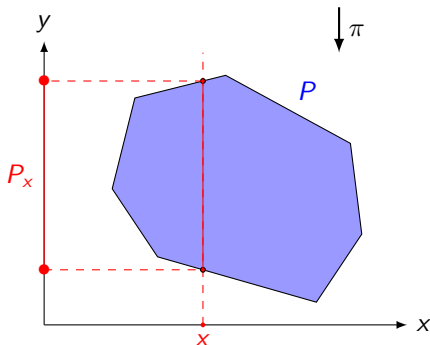
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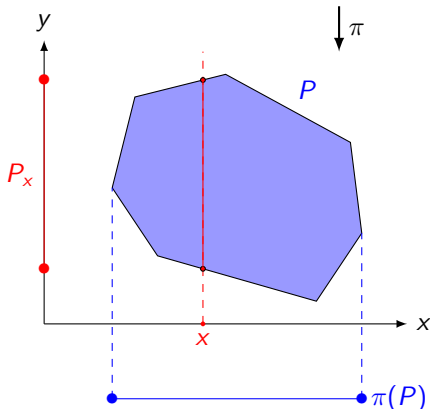
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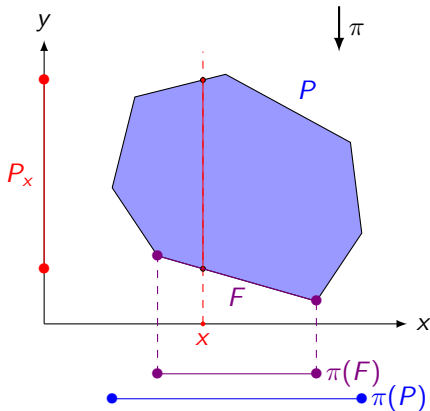
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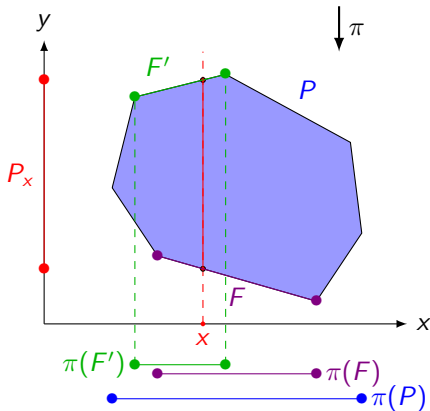
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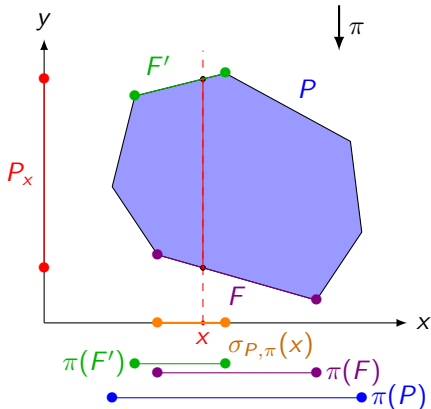
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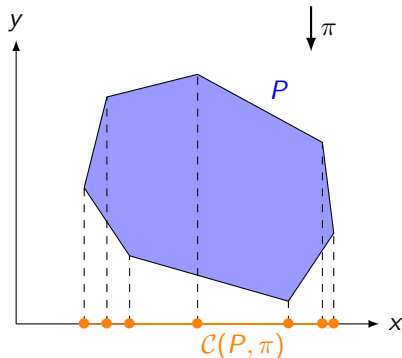
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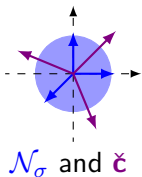
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General cost \mathbf{c} is equivalent to discrete cost $\check{\mathbf{c}}$ for all x

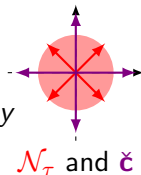


For all $x \in \text{ri}(\sigma)$,

$$V(x) = \sum_{N \in \mathcal{N}_\sigma} p_N \min_{y \in P_x} \check{\mathbf{c}}_N^\top y$$

For all $x \in \text{ri}(\tau)$,

$$V(x) = \sum_{N \in \mathcal{N}_\tau} p_N \min_{y \in P_x} \check{\mathbf{c}}_N^\top y$$



Theorem (Quantization of the cost distribution)

Let $\mathcal{R} = \bigwedge_{\sigma \in \mathcal{C}(P, \pi)} -\mathcal{N}_\sigma$, then for all $x \in \mathbb{R}^n$

$$V(x) = \sum_{R \in \mathcal{R}} \check{p}_R \min_{y \in \mathbb{R}^m} \check{\mathbf{c}}_R^\top y + \mathbb{I}_{T_x + W y \leq h}$$

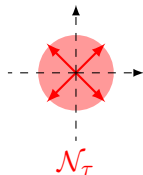
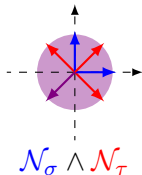
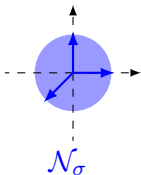
where $\check{p}_R := \mathbb{P}[\mathbf{c} \in \text{ri}(R)]$ and $\check{\mathbf{c}}_R := \mathbb{E}[\mathbf{c} \mid \mathbf{c} \in \text{ri}(R)]$

Bonus: This quantization method works for *every distribution of \mathbf{c}* !

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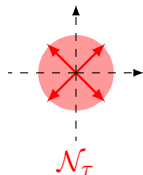
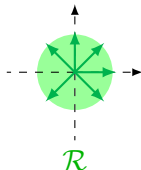
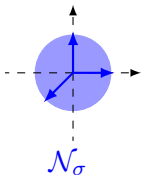
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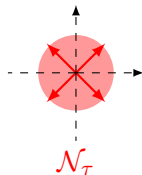
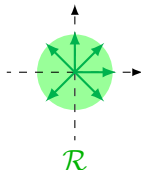
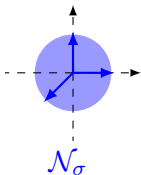
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Exact formula of $V(x)$ for all x

There exists a collection \mathcal{I}_σ of active constraints sets such that $\mathcal{N}_\sigma = \{\text{Cone}(W_I^\top) \mid I \in \mathcal{I}_\sigma\}$.

Theorem

For all $I \in \mathcal{I}_\sigma$, there exists $\mu(I) \in \mathbb{R}_+^I$ such that

$$-\mu(I)^\top W_I = \mathbb{E}[\mathbf{1}_{\mathbf{c} \in -\text{ri Cone}(W_I^\top)} \mathbf{c}^\top]$$

Let $\alpha_\sigma := \sum_{I \in \mathcal{I}_\sigma} T_I^\top \mu(I)$ and $\beta_\sigma := -\sum_{I \in \mathcal{I}_\sigma} h_I^\top \mu(I)$. Then, for all $x \in \mathbb{R}^n$,

$$V(x) = \max_{\sigma \in \mathcal{C}^{\max}(P, \pi)} (\alpha_\sigma^\top x + \beta_\sigma) + \mathbb{I}_{x \in \pi(P)}$$

In particular, V is polyhedral.

Bonus: for all distributions of \mathbf{c} , V is affine on each cell of $\mathcal{C}(P, \pi)$.

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Extension to multistage and stochastic constraints

Theorem

All results generalize to stochastic constraints with finite support and multistage

- \rightsquigarrow *All V_t are polyhedral* (easy)
- \rightsquigarrow *The regions where $(V_t)_t$ is affine do not depend on the $(c_t)_t$ (harder)*
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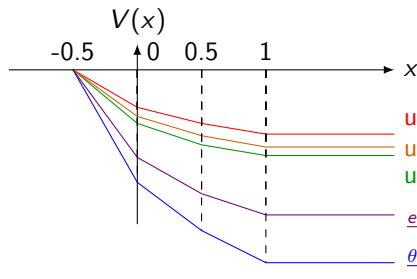
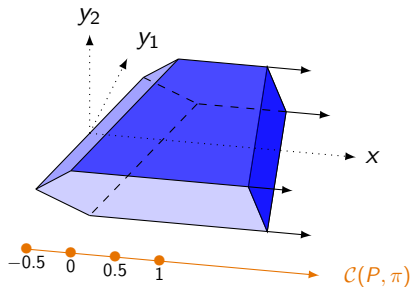
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Explicit computation of the example

$$V(x) = \mathbb{E} \left[\begin{array}{ll} \min_{y \in \mathbb{R}^2} & \mathbf{c}^\top y \\ \text{s.t.} & \|y\|_1 \leq 1 \\ & y_1 \leq x \\ & y_2 \leq x \end{array} \right]$$



Different distributions of \mathbf{c} :

uniform on norm 1 ball

uniform on norm 2 ball

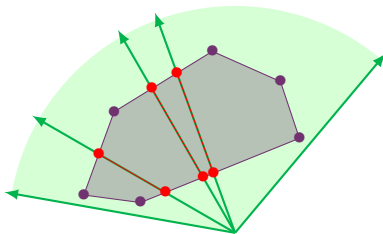
uniform on norm ∞ ball

$$\frac{e^{-\frac{\|\mathbf{c}\|_2^2}{2\gamma^2}}}{\theta^2 e^{-\theta \|\mathbf{c}\|_1}} d\mathbf{c}$$

We can triangulate $-N \cap \text{supp}(\mathbf{c})$

We need to compute $\mathbb{E}[\mathbf{c} \mathbf{1}_{\mathbf{c} \in \text{ri}(N)}] = \mathbb{P}[\mathbf{c} \in \text{ri}(N)] \mathbb{E}[\mathbf{c} \mid \mathbf{c} \in \text{ri}(N)]$ which is additive.

The shape of $-N \cap \text{supp}(\mathbf{c})$ can be complicated.



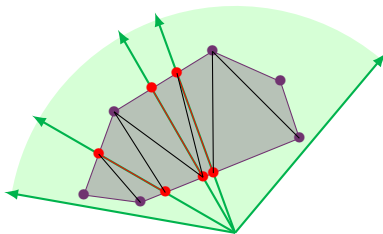
$$\mathcal{N}(P_x) \cap -\text{supp}(\mathbf{c})$$

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\rightsquigarrow We can triangulate



$$\mathcal{N}(P_x) \cap -\text{supp}(\mathbf{c})$$

Explicit formulas for usual distributions

We can compute explicitly $\mathbb{E}[\mathbf{c} \mathbb{1}_{\mathbf{c} \in \text{ri}(N)}] = \mathbb{P}[\mathbf{c} \in \text{ri}(N)] \mathbb{E}[\mathbf{c} \mid \mathbf{c} \in \text{ri}(N)]$ for classical distributions

Distribution	Uniform on polytope	Exponential	Gaussian
$d\mathbb{P}(\mathbf{c})$	$\frac{\mathbb{1}_{\mathbf{c} \in Q}}{\text{Vol}_d(Q)} d\mathcal{L}_{\text{Aff}(Q)}(\mathbf{c})$	$\frac{e^{\theta^\top \mathbf{c}} \mathbb{1}_{\mathbf{c} \in K}}{\Phi_K(\theta)} d\mathcal{L}_{\text{Aff}(K)}(\mathbf{c})$	$\frac{e^{-\frac{1}{2} \mathbf{c}^\top M^{-2} \mathbf{c}}}{(2\pi)^{\frac{m}{2}} \det M} d\mathbf{c}$
Support	Polytope : Q	Cone : K	\mathbb{R}^m
$\mathbb{P}[\mathbf{c} \in S]$	$\frac{\text{Vol}_d(S)}{\text{Vol}_d(Q)}$	$\frac{ \det(\text{Ray}(S)) }{\Phi_K(\theta)} \prod_{r \in \text{Ray}(S)} \frac{1}{-r^\top \theta}$	$\text{Ang}(M^{-1}S)$
$\mathbb{E}[\mathbf{c} \mid \mathbf{c} \in S]$	$\frac{1}{d} \sum_{v \in \text{Vert}(S)} v$	$\left(\sum_{r \in \text{Ray}(S)} \frac{-r_i}{r^\top \theta} \right)_{i \in [m]}$	$\frac{\sqrt{2} \Gamma(\frac{m+1}{2})}{\Gamma(\frac{m}{2})} M \text{Centr}(S \cap \mathbb{S}_{m-1})$

These formulas are valid for S full dimensional **simplex** or **simplicial cone**.

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Earlier and new complexity results

Volume of a polytope

$$\text{Vol}(\{z \in \mathbb{R}^d \mid Az \leq b\}) \text{ or} \\ \text{Vol}(\text{Conv}(v_1, \dots, v_n))$$

- $\#P$ -complete: Dyer and Frieze (1988)
- Polynomial for fixed dimension d : Barvinok (1994)

2-stage linear problem

$$\min_{x \in \mathbb{R}^n} c_0^\top x + \mathbb{I}_{Ax \leq b} \\ + \mathbb{E} \left[\min_{y \in \mathbb{R}^m} c^\top y + \mathbb{I}_{Tx + Wy \leq h} \right]$$

- $\#P$ -hard: Hanasusanto, Kuhn and Wiese (2016)
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- $\#P$ -hard: Hanasusanto, Kuhn and Wiese (2016)
- **Polynomial for fixed m** for some usual distributions: FGL (2020)

Complexity results

We make the following assumption:

$\#\text{supp}(\mathbf{T}, \mathbf{W}, \mathbf{h})$ is finite and $\mathbb{P}[\mathbf{c} \in S]$ and $\mathbb{E}[\mathbf{c} \mid \mathbf{c} \in S]$ can be computed in polynomial time for S full dimensional simplex or simplicial fan.

Theorem

$$\min_{x \in \mathbb{R}^n} c_0^\top x + \mathbb{I}_{Ax \leq b} + \mathbb{E} \left[\min_{y \in \mathbb{R}^m} \mathbf{c}^\top y + \mathbb{I}_{\mathbf{T}x + \mathbf{W}y \leq \mathbf{h}} \right] \quad (2\text{SLP})$$

When m is fixed, we can solve (2SLP) in polynomial time.

Theorem

For $x \in \mathbb{R}^n$

$$V(x) := \mathbb{E} \left[\min_{y \in \mathbb{R}^m} \mathbf{c}^\top y + \mathbb{I}_{\mathbf{T}x + \mathbf{W}y \leq \mathbf{h}} \right]$$

When n , m and $\#\text{supp}(\mathbf{T}, \mathbf{W}, \mathbf{h})$ are fixed, we can compute the epigraph of V in polynomial time.

Complexity result multistage

We can generalize to multistage by fixing several dimensions and the horizon.

Theorem (MSLP is polynomial for fixed dimensions)

*Assume that $t_{\max} \geq 3$, $n_2, \dots, n_{t_{\max}}$, $\sharp(\text{supp}(\mathbf{T}_2, \mathbf{W}_2, \mathbf{h}_2))$,
 $\dots, \sharp(\text{supp}(\mathbf{T}_{t_{\max}}, \mathbf{W}_{t_{\max}}, \mathbf{h}_{t_{\max}}))$ are fixed integers*

and for all $t \in [t_{\max}]$, \mathbf{c}_t conditionally to $\{(\mathbf{T}_t, \mathbf{W}_t, \mathbf{h}_t) = (T, W, h)\}$ is easily computable.

Then, we can solve MSLP in polynomial time.

Conclusion

- V_t are polyhedral
and affine on regions that are independent of the cost distribution;
- MSLP with arbitrary cost distribution can be exactly discretized;
- analytical formulas for some usual distributions;
- fixed-parameter versions of 2SLP and MSLP are polynomial time.

Perspectives

- ↪ Deduce a faster algorithm from the algebraic structure
- ↪ Extend the complexity analysis to approximation of the problem;
- ↪ Extend to integer stochastic problems;
- ↪ Distributionnally robust optimization or sensibility analysis.

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References

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Thank you for listening ! Any question ?

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