Multistage Stochastic Linear Problem and Polyhedral Geometry

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Multistage stochastic linear programming (MSLP)

$$\min_{(\mathbf{x}_t)_{t \in [T]}} \mathbb{E} \Big[\sum_{t=1}^{T} \mathbf{c}_t^\top \mathbf{x}_t \Big]$$
s.t. $\mathbf{A}_t \mathbf{x}_t + \mathbf{B}_t \mathbf{x}_{t-1} \leq \mathbf{b}_t \qquad \forall t \in [T]$
 \mathbf{x}_t random variable in $\mathbb{R}^{n_t} \qquad \forall t \in [T]$
 $\sigma(\mathbf{x}_t) \subset \sigma(\mathbf{c}_k, \mathbf{A}_k, \mathbf{B}_k, \mathbf{b}_k)_{k \leq t} \qquad \forall t \in [T]$
 $\mathbf{x}_0 \equiv \mathbf{x}_0$ given

where $\mathbf{c}_t \in \mathbb{R}^{n_t}$, $\mathbf{A}_t \in \mathbb{R}^{q_t \times n_{t-1}}$, $\mathbf{B}_t \in \mathbb{R}^{q_t \times n_t}$ and $\mathbf{b}_t \in \mathbb{R}^{q_t}$ are given random variables.

 $(\mathbf{c}_t, \mathbf{A}_t, \mathbf{B}_t, \mathbf{b}_t)_{t \in [T]}$ is an independent sequence.

We set $V_{T+1} \equiv 0$ and:

$$V_t(x_{t-1}) := \mathbb{E} \begin{bmatrix} \min_{x_t \in \mathbb{R}^{n_t}} \mathbf{c}_t^\top x_t + V_{t+1}(x_t) \\ \text{s.t. } \mathbf{A}_t x_t + \mathbf{B}_t x_{t-1} \leqslant \mathbf{b}_t \end{bmatrix}$$

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Quantization of a MSLP

The random variable $(\mathbf{c}_t, \mathbf{A}_t, \mathbf{B}_t, \mathbf{b}_t)_{t \in [T]}$ are often replaced by a discrete distribution on a finite number of scenarios

$$V_t(x_{t-1})\simeq \widetilde{V}_t(x_{t-1}) = \sum_{k=1}^K p_k \sum_{x_t\in\mathbb{R}^{n_t}}^{\min} c_{t,k}^ op x_t + V_{t+1}(x_t) \ ext{s.t. } A_{t,k}x_t + B_{t,k}x_{t-1}\leqslant b_{t,k}$$

Scenario drawn by Monte Carlo : Sample Average Approximation

Definition

We say that an MSLP admits an *exact quantization* if there exists a finitely supported $(\check{\mathbf{c}}_t, \check{\mathbf{A}}_t, \check{\mathbf{B}}_t, \check{\mathbf{b}}_t)_{t \in [T]}$ that yields the same expected cost-to-go functions, $(V_t)_{t \in [T]}$. In particular the MSLP is equivalent to a problem on a finite scenario tree.

Contents

Exact Quantization Result

- Fixed state x and normal fan
- Variable state x and chamber complex

2 Fiber Polyhedron

3 Complexity results

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$$V(x) = \mathbb{E} \begin{bmatrix} \min_{y \in \mathbf{P}_x} \mathbf{c}^\top y \end{bmatrix} \text{ where } \mathbf{P}_x := \{ y \in \mathbb{R}^m \mid Ax + By \leq b \}$$

Illustrative running example:



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Definition

The normal fan of the fiber P_x is

$$\mathcal{N}(P_x) := \{N_{P_x}(y) \mid y \in P_x\}$$



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 y and $N_{P_x}(y)$ for $x = 0.3$

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with $N_{P_x}(y) = \{c \mid \forall y' \in P_x, c^{\top}(y'-y) \leq 0\}$ the normal cone of P_x on y.



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with $N_{P_x}(y) = \{ c \mid \forall y' \in P_x, \ c^\top(y'-y) \leqslant 0 \}$ the normal cone of P_x on y.

Proposition

If P_x is bounded, $\{ri(N) | N \in \mathcal{N}(P_x)\}$ is a partition of \mathbb{R}^m .





$$V(x) = \mathbb{E}\big[\min_{y\in P_x} \mathbf{c}^{\top}y\big]$$

For any $N \in \mathcal{N}(P_x)$ and $-c \to \arg \min_{y \in P_x} c^\top y$ is constant for all $-c \in \operatorname{ri}(N)$.

 $\arg\min_{y\in P_x} c^{\top} y$ is a face of P_x .





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Reduction to a finite sum

For a fixed x,

$$V(x) = \mathbb{E}\big[\min_{y \in P_x} \mathbf{c}^\top y\big] = \sum_{N \in \mathcal{N}(P_x)} \mathbb{E}\big[\mathbf{c}^\top \mathbb{1}_{\mathbf{c} \in -\operatorname{ri} N}\big] y_N(x)$$

where $y_N(x) \in \arg \min_{y \in P_x} c^\top y$ for any $c \in ri(N)$.





General cost **c** is equivalent to discrete cost $\check{\mathbf{c}}$ for given xFor a fixed x,

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$$= \sum_{N \in \mathcal{N}(P_{x})} \mathbb{E}\left[\mathbf{c}^{\top} \mathbb{1}_{\mathbf{c} \in -\operatorname{ri} N}\right] y_{N}(x)$$



 $\mathcal{N}(P_x)$ for x = 0.3

We draw a continuous cost **c**.

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$$= \sum_{N \in \mathcal{N}(P_{x})} p_{N} \check{c}_{N}^{\top} y_{N}(x)$$



 $\mathcal{N}(P_x)$ and $p_N \check{c}_N$ for x = 0.3

For $N \in \mathcal{N}(P_x)$,

$$p_N := \mathbb{P}\big[\mathbf{c} \in -\operatorname{ri} N\big]$$
$$\check{c}_N := \mathbb{E}\big[\mathbf{c} | \mathbf{c} \in -\operatorname{ri} N\big]$$

Instead of drawing a general \mathbf{c} , we draw a discrete cost $\check{\mathbf{c}}$ indexed by the finite collection $\mathcal{N}(P_{\times})$. General cost **c** is equivalent to discrete cost $\check{\mathbf{c}}$ for given x For a fixed x,

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= $\sum_{N \in \mathcal{N}(P_{x})} p_{N} \min_{y \in P_{x}} \check{c}_{N}^{\top} y$

For $N \in \mathcal{N}(P_x)$,

$$p_N := \mathbb{P} \big[\mathbf{c} \in -\operatorname{ri} N \big] \\ \check{c}_N := \mathbb{E} \big[\mathbf{c} | \mathbf{c} \in -\operatorname{ri} N \big]$$





Instead of drawing a general \mathbf{c} , we draw a discrete cost $\check{\mathbf{c}}$ indexed by the finite collection $\mathcal{N}(P_{\times})$.
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 $P := \{(x, y) \mid Ax + By \leq b\} \text{ and } P_x := \{y \mid Ax + By \leq b\}$ x = -0.1*Y*2 *Y*2 *Y*1 $-c_{2}$ $-c_1$ ► X P_x and $\mathcal{N}(P_x)$ x = -0.1 $\mathcal{N}(P_x)$

 $P := \{(x, y) \mid Ax + By \leq b\}$ and $P_x := \{y \mid Ax + By \leq b\}$ x = 0*y*₂ *Y*2 *Y*1 $-c_{2}$ ► *Y*1 c_1 ► X $\mathcal{N}(P_x)$ P_x and $\mathcal{N}(P_x)$ x = 0

 $P := \{(x, y) \mid Ax + By \leq b\}$ and $P_x := \{y \mid Ax + By \leq b\}$ x = 0.1 y_2 *Y*2 *Y*1 $-c_{2}$ ► *Y*1 C_1 ► X $\mathcal{N}(P_x)$ P_x and $\mathcal{N}(P_x)$ x = 0.1



 $P := \{(x, y) \mid Ax + By \leq b\} \text{ and } P_x := \{y \mid Ax + By \leq b\}$ x = 0.2 y_2 *Y*2 *Y*1 $-c_{2}$ *Y*1 c_1 ► X $\mathcal{N}(P_x)$ P_x and $\mathcal{N}(P_x)$ *x* = 0.2

P and P_x

 $P := \{(x, y) \mid Ax + By \leq b\} \text{ and } P_x := \{y \mid Ax + By \leq b\}$ x = 0.3 y_2 *Y*2 У1 $-c_{2}$ y_1 ► X $\mathcal{N}(P_x)$ P_x and $\mathcal{N}(P_x)$ *x* = 0.3



 $P := \{(x, y) \mid Ax + By \leq b\}$ and $P_x := \{y \mid Ax + By \leq b\}$ x = 0.4 y_2 *Y*2 y_1 $-c_{2}$ <u>--</u>- + *Y*1 ► X $\mathcal{N}(P_x)$ P_x and $\mathcal{N}(P_x)$ x = 0.4

 $P := \{(x, y) \mid Ax + By \leq b\}$ and $P_x := \{y \mid Ax + By \leq b\}$ x = 0.5 y_2 y_1 $-c_{2}$ y_1 ► X $\mathcal{N}(P_x)$ P_x and $\mathcal{N}(P_x)$ x = 0.5

P and P_x

 $P := \{(x, y) \mid Ax + By \leq b\}$ and $P_x := \{y \mid Ax + By \leq b\}$ x = 0.6 y_2 y_1 $-c_{2}$ $\rightarrow y_1$ c_1 ► X $\mathcal{N}(P_x)$ P_x and $\mathcal{N}(P_x)$ *x* = 0.6

 $P := \{(x, y) \mid Ax + By \leq b\}$ and $P_x := \{y \mid Ax + By \leq b\}$ x = 0.7*Y*2 y_1 $-c_{2}$ y_1 c_1 ► X $\mathcal{N}(P_x)$ P_x and $\mathcal{N}(P_x)$ x = 0.7



 $P := \{(x, y) \mid Ax + By \leq b\}$ and $P_x := \{y \mid Ax + By \leq b\}$ x = 0.9*Y*2 *Y*1 $-c_{2}$ · C1 ► X $\mathcal{N}(P_x)$ P_x and $\mathcal{N}(P_x)$ x = 0.9**P** and P_{x}





 $P := \{(x, y) \mid Ax + By \leq b\}$ and $P_x := \{y \mid Ax + By \leq b\}$ x = 1.2 y_2 *Y*1 $-c_{2}$ с1 ► X $\mathcal{N}(P_x)$ P_x and $\mathcal{N}(P_x)$ x = 1.2**P** and P_{x}

 $P := \{(x, y) \mid Ax + By \leq b\}$ and $P_x := \{y \mid Ax + By \leq b\}$ x = 1.3 y_2 *Y*1 $-c_{2}$ *c*₁ ► X $\mathcal{N}(P_x)$ P_x and $\mathcal{N}(P_x)$ x = 1.3**P** and P_{x}

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P and P_x

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What are the constant regions of $\mathcal{N}(P_x)$?

Lemma

There exists a collection $C(P, \pi)$ called the **chamber complex** whose relative interior of cells are the constant regions of $x \to \mathcal{N}(P_x)$.

For $\sigma \in C(P, \pi)$ and $x, x' \in ri(\sigma)$, $\mathcal{N}(P_x) = \mathcal{N}(P_{x'}) =: \mathcal{N}_{\sigma}$





Definition (Billera, Sturmfels 92) The chamber complex $C(P, \pi)$ of Palong π is $C(P, \pi) := \{\sigma_{P,\pi}(x) \mid x \in \pi(P)\}$ where $\sigma_{P,\pi}(x) := \bigcap_{F \in \mathcal{F}(P) \text{ s.t. } x \in \pi(F)} \pi(F)$



$$\pi(E) := \{x \in \mathbb{R}^n \mid \exists y \in \mathbb{R}^m, \ (x, y) \in E\}$$



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where $\mathcal{F}(P)$ is the set of faces of Pand π is the projection $(x, y) \rightarrow x$

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 π

1.1

Common Refinement of Normal Fans

We can quantize **c** on each chamber.

For all
$$x \in \operatorname{ri}(\sigma)$$
, For all $x \in \operatorname{ri}(\tau)$,
 $V(x) = \sum_{N \in \mathcal{N}_{\sigma}} p_{N} \min_{y \in P_{x}} \check{c}_{N}^{\top} y$ $V(x) = \sum_{N \in \mathcal{N}_{\tau}} p_{N} \min_{y \in P_{x}} \check{c}_{N}^{\top} y$
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We take the *common refinement*:

$$\mathcal{R} := \mathcal{N}_{\sigma} \land \mathcal{N}_{\tau} = \{ \mathcal{N} \cap \mathcal{N}' \mid \mathcal{N} \in \mathcal{N}_{\sigma}, \mathcal{N}' \in \mathcal{N}_{\tau} \}$$



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General cost \mathbf{c} is equivalent to discrete cost $\check{\mathbf{c}}$ for all x

Theorem (Quantization of the cost distribution) Let $\mathcal{R} = \bigwedge_{\sigma \in \mathcal{C}(P,\pi)} -\mathcal{N}_{\sigma}$, then for all $x \in \mathbb{R}^{n}$ $V(x) = \sum_{R \in \mathcal{R}} \check{p}_{R} \min_{y \in P_{x}} \check{c}_{R}^{\top} y$ where $\check{p}_{R} := \mathbb{P}[\mathbf{c} \in \operatorname{ri}(R)]$ and $\check{c}_{R} := \mathbb{E}[\mathbf{c} | \mathbf{c} \in \operatorname{ri}(R)]$ Moreover, for all distributions of \mathbf{c} , V is affine on each cell of the chamber complex $\mathcal{C}(P,\pi)$.

Bonus: This quantization method works for *every distribution of* **c** !

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Extension to multistage and stochastic constraints

Theorem

All results generalize to stochastic constraints with finite support and multistage

- \rightsquigarrow The regions where $(V_t)_t$ is affine do not depend on the $(\boldsymbol{c}_t)_t$
- \rightsquigarrow We have an exact discretization method working for all $(c_t)_t$

Idea of the proof : Iterated chamber complexes

$$\mathcal{P}_{t,\xi} := \mathcal{C}(\mathbb{R}^{n_t} \times \mathcal{P}_{t+1} \wedge \mathcal{F}(P_t(\xi)), \pi_{x_{t-1}}^{x_{t-1}, x_t})$$
$$\mathcal{P}_t := \bigwedge_{\xi_t \in \text{supp } \xi_t} \mathcal{P}_{t,\xi}$$

Explicit computation of the example

$$V(x) = \mathbb{E} \begin{bmatrix} \min_{y \in \mathbb{R}^2} & \mathbf{c}^\top y \\ \text{s.t. } \|y\|_1 \leq 1 \\ y_1 \leq x \\ y_2 \leq x \end{bmatrix}$$





Explicit formulas for usual distributions

Distribution	Uniform on polytope	Exponential	Gaussian
$d\mathbb{P}(c)$	$rac{\mathbb{1}_{c\in Q}}{\operatorname{Vol}_d(Q)} d\mathcal{L}_{\operatorname{Aff}(Q)}(c)$	$\frac{e^{\theta^{\top}c}\mathbb{1}_{c\in K}}{\Phi_{K}(\theta)}d\mathcal{L}_{\mathrm{Aff}(K)}c$	$\frac{e^{-\frac{1}{2}c^{\top}M^{-2}c}}{(2\pi)^{\frac{m}{2}}\det M}dc$
Support	Polytope : Q	Cone : K	\mathbb{R}^{m}
$\mathbb{P}ig[\mathbf{c}\in Sig]$	$\frac{\operatorname{Vol}_d(S)}{\operatorname{Vol}_d(Q)}$	$\frac{ \det(Ray(S)) }{\Phi_{\mathcal{K}}(\theta)} \prod_{r \in Ray(S)} \frac{1}{-r^{\top}\theta}$	$\operatorname{Ang}\left(M^{-1}S\right)$
$\mathbb{E}\left[\mathbf{c} \mid \mathbf{c} \in S ight]$	$\frac{1}{d}\sum_{v\in \operatorname{Vert}(S)} v$	$\left(\sum_{r\inRay(S)}\frac{-r_i}{r^{\top}\theta}\right)_{i\in[m]}$	$\frac{\sqrt{2}\Gamma(\frac{m+1}{2})}{\Gamma(\frac{m}{2})}M\operatorname{Ctr}\left(S\cap\mathbb{S}_{m-1}\right)$

These formulas are valid for S full dimensional simplex or simplicial cone.

Contents

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2 Fiber Polyhedron

B Complexity results

Dual problem

$$V(x) := \mathbb{E} \begin{bmatrix} \inf_{y} & \mathbf{c}^{\top} y \\ s.t. & Ax + By \leq b \end{bmatrix} = \mathbb{E} \begin{bmatrix} \inf_{y \in P_{x}} \mathbf{c}^{\top} y \end{bmatrix}$$

where $P_x = \{x \mid Ax + By \leq b\}$

$$V(x) := \mathbb{E} \begin{bmatrix} \sup_{\mu} & (Ax - b)^{\top} \mu \\ \text{s.t.} & B^{\top} \mu + \mathbf{c} = 0 \\ & \mu \ge 0 \end{bmatrix} = \mathbb{E} \begin{bmatrix} \sup_{\mu \in D_{\mathbf{c}}} (Ax - b)^{\top} \mu \end{bmatrix}$$

where $D_c = \{\mu \mid B^\top \mu + c = 0, \mu \ge 0\}$

Minkowski sum :

$$E + F = \{x + x' | x \in E, x' \in F\}$$

Definition

$$\boldsymbol{E} := \int \boldsymbol{D}_{\boldsymbol{c}} \mathbb{P}(\boldsymbol{d}\boldsymbol{c}) = \left\{ \int \mu(\boldsymbol{c}) \mathbb{P}(\boldsymbol{d}\boldsymbol{c}) \, | \, \mu(\boldsymbol{c}) \in \boldsymbol{D}_{\boldsymbol{c}} \quad \boldsymbol{a.s.}, \ \mu \in L_{\infty}(\mathbb{R}^{m}, \mathbb{R}^{l}) \right\}$$

$$\begin{split} \ell(x) &= \mathbb{E} \Big[\sup_{\mu \in D_{\mathbf{c}}} (Ax - b)^{\top} \mu \Big] \\ &= \begin{cases} \sup_{\mu(\cdot)} & (Ax - b)^{\top} \mathbb{E} \big[\mu(\mathbf{c}) \big] \\ \text{s.t.} & \mu(\mathbf{c}) \in D_{\mathbf{c}} \text{ a.s.} \end{cases} \end{split}$$

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The Fiber Polyhedron is a finite Minkowski sum

Theorem

There exists a chamber complex \mathcal{R} depending on A such that

$$\boldsymbol{E} = \int \boldsymbol{D}_{\boldsymbol{c}} \mathbb{P}(\boldsymbol{d}\boldsymbol{c}) = \sum_{\boldsymbol{R} \in \mathcal{R}} \check{p}_{\boldsymbol{R}} \boldsymbol{D}_{\check{\boldsymbol{c}}_{\boldsymbol{R}}}$$

where $\check{p}_R := \mathbb{P} \big[\mathbf{c} \in \mathsf{ri}(R) \big]$ and $\check{c}_R := \mathbb{E} \big[\mathbf{c} \,|\, \mathbf{c} \in \mathsf{ri}(R) \big]$.

Alternative proof of the quantization result

$$V(x) = \sigma_E(Ax - b) = \sum_{R \in \mathcal{R}} \check{p}_R \sigma_{D_{\check{c}_R}}(Ax - b) = \sum_{R \in \mathcal{R}} \check{p}_R \min_{y \in P_x} \check{c}_R^\top y$$

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Nested Fiber Polyhedra for Multistage

$$V_t(x_{t-1}) = \mathbb{E} \begin{bmatrix} \min_{x_t \in \mathbb{R}^{n_t}} \mathbf{c}_t^\top x_t + V_{t+1}(x_t) \\ \text{s.t. } A_t x_t + B_t x_{t-1} \leqslant b_t \end{bmatrix}$$

Definition

We define by induction the following nested fiber polyhedra

$$\begin{aligned} D_{t,c_{t}} &:= \{\mu_{t} | \mu_{t} \ge 0, A_{t}^{\top} \mu_{t} + c_{t} = 0\} & \forall t \in [T] \\ F_{T,c_{T}} &:= D_{T,c_{T}} \\ E_{t} &:= \mathbb{E} [F_{t,c_{t}}] & \forall t \in [T] \\ F_{t,c_{t}} &:= \{(\mu_{t}, \lambda_{[t+1:T]}) | \mu_{t} \in D_{t,c_{t}} + B_{t+1}^{\top} \lambda_{t+1}, \ \lambda_{[t+1:T]} \in E_{t+1}\} & \forall t \in [T-1] \end{aligned}$$

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2-time scale MSLP reduced to Quadratic Problem

We consider a MSLP where the dynamic depends on parameters p we have to optimize

$$\min_{p \in \mathbb{R}^{m}, (\mathbf{x}_{t}) \in \mathbb{R}^{n_{t}}} \quad q^{\top} p + \mathbb{E} \Big[\sum_{t=1}^{T} \mathbf{c}_{t}^{\top} \mathbf{x}_{t} \Big]$$
s.t. $Dp \leq d$
 $A_{t} \mathbf{x}_{t} + B_{t} \mathbf{x}_{t-1} + C_{t} p \leq h_{t}$ a.s. $\forall t \in [T]$
 $\mathbf{x}_{t} \prec \sigma(\mathbf{c}_{1}, \cdots, \mathbf{c}_{t}) \qquad \forall t \in [T]$

If we know the fiber polyhedron, it reduces to a finite dimensional quadratic problem

$$\sup_{\substack{p \in \mathbb{R}^{m}, (\lambda_{t})_{t \in [T]}}} -q^{\top}p + \sum_{t=1}^{T} (C_{t}p - h_{t})^{\top}\lambda_{t}$$

s.t. $Dp \leq d$
 $(\lambda_{1}, \cdots, \lambda_{T}) \in E_{1}$

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Earlier and new complexity results

Volume of a polytope

$$\mathsf{Vol}\left(\{z\in\mathbb{R}^d\,|\,Az\leqslant b\}
ight)$$
 or $\mathsf{Vol}\left(\mathsf{Conv}(v_1,\cdots,v_n)
ight)$

- #*P*-complete: Dyer and Frieze (1988)
- Polynomial for fixed dimension *d*: Barvinok (1994)

$$\min_{x \in \mathbb{R}^n} c_0^\top x + \mathbb{I}_{Ax \leqslant b} \\ + \mathbb{E} \big[\min_{y \in \mathbb{R}^m} \mathbf{c}^\top y + \mathbb{I}_{\mathbf{T}x + \mathbf{W}y \leqslant \mathbf{h}} \big]$$

- *#P*-hard: Hanasusanto, Kuhn and Wiesemann (2016)
- Polynomial for fixed *m* ?

Earlier and new complexity results

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- *#P*-hard: Hanasusanto, Kuhn and Wiesemann (2016)
- Polynomial for fixed *m*: FGL (2020)

Complexity result multistage

We can generalize to multistage by fixing several dimensions and the horizon.

Theorem (MSLP is polynomial for fixed dimensions)

Assume that $T \ge 3, n_2, \ldots, n_T$, $\sharp (supp(\mathbf{A}_2, \mathbf{B}_2, \mathbf{b}_2)), \cdots, \sharp (supp(\mathbf{A}_T, \mathbf{B}_T, \mathbf{b}_T))$ are fixed integers

and for all $t \in [T]$, \mathbf{c}_t conditionally to $\{(\mathbf{A}_t, \mathbf{B}_t, \mathbf{b}_t) = (A, B, b)\}$ is easily computable.

Then, we can solve MSLP in polynomial time.

Conclusion

• MSLP with arbitrary cost distribution can be exactly discretized;

- new algebraic insights on the polyhedral structure of MSLP;
- analytical formulas for some usual distributions;
- fixed-parameter versions of 2SLP and MSLP are polynomial time.

Perspectives

- \rightsquigarrow New algorithms from the algebraic structure
- \rightsquigarrow Sensibility analysis to the distribution, link with nested distance;
- → Extend to integer stochastic problems;
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Thank you for listening ! Any question ?

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